# **ER-to-Relational Mapping**

#### Lecture 8



**ER-to-Relational Mapping** 

March 1, 2018

#### Outline

- 1. Context
- 2. The Algorithm



#### **Database Design and Implementation Process**





#### Data Models









# **Resulting Relational Schema**

#### **EMPLOYEE**





## Step 1: Regular Entity Types

- i. For each regular/strong entity type, create a corresponding relation that includes all the <u>simple</u> attributes (includes simple attributes of composite relations)
- ii. Choose one of the key attributes as primary
  - If composite, the simple attributes together form the primary key
- iii. Any remaining key attributes are kept as secondary unique keys (these will be useful for physical tuning w.r.t. indexing analysis)







## Step 1 Result

#### EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary
DEPARTM	IENT						
Dname	Dnumb	<u>er</u>					
PROJECT							
Pname	Pnumb	er Ploc	cation				



# Step 2: Weak Entity Types

- i. For each weak entity type, create a corresponding relation that includes all the simple attributes
- ii. Add as a foreign key all of the primary key attribute(s) in the entity corresponding to the owner entity type
- iii. The primary key is the combination of all the primary key attributes from the owner and the partial key of the weak entity, if any







## Step 2 Result

#### **EMPLOYEE**

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	
DEPARTMENT								
Dname	Dnumb	<u>per</u>						
PROJECT	-							
Pname Pnumber Plocation								
DEPENDENT								
Essn	Dependent_name		Sex	Bdate	Relations	ship		



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# Step 3: Mapping Binary 1-to-1

#### Three approaches

#### – Foreign Key

- Usually appropriate
- Merged Relation
  - Possible when both participations are total
- Relationship Relation
  - Not discussed



#### Step 3: Mapping Binary 1-to-1 Foreign Key

- i. Choose one relation as S, the other T
  - Better if S has total participation (reduces number of NULL values)
- ii. Add to S all the simple attributes of the relationship
- iii. Add as a foreign key in S the primary key attributes of T







### Step 2 Result

#### EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary
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#### DEPARTMENT

Dname <u>Dnumber</u>



### Step 3 Result

EMPLOYEE

Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary
DEPARTMENT							
Dname	Dnumb	<u>er</u> Mgı	_ssn	Mgr_start	_date		



## Step 4: Binary 1-to-N

i. Choose the S relation as the type at the N-side of the relationship, other is T

ii. Add as a <u>foreign key</u> to S all of the primary key attribute(s) of T

Another approach: create a relationship relation







### Step 4 Result

EMF	PLOY	ΈΕ
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# Step 5: Binary M-to-N

- i. Create a <u>new</u> relation S (termed: *relationship relation*)
  - In some ERD dialects, actually drawn in
- ii. Add as foreign keys the primary keys of both relations; their <u>combination</u> forms the primary key of S
- iii. Add any simple attributes of the M:N relationship to S







#### Step 5 Result







## Step 6: Multivalued Attributes

- i. Create a <u>new</u> relation S
- ii. Add as foreign keys the primary keys of the corresponding relation

iii. Add the attribute to S (if composite, the simple attributes); the combination of all attributes in S forms the primary key







## Step 6 Result







#### Step 7: Specialization/Generalization

- A. Multiple relations subclass and superclass
  - Usually works (assumes unique id at parent)
- B. Multiple relations subclass only
  - Should only be used for disjoint
- C. Single relation with one type attribute
  - Only for disjoint, can result in many NULLs
- D. Single relation with multiple type attributes
  - Better for overlapping, could be disjoint



# Specialization/Generalization (A)





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## Specialization/Generalization (B)





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# Specialization/Generalization (C)





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## Specialization/Generalization (D)

PART

Part\_no Description Mflag Drawing\_no Manufacture\_date Batch\_no Pflag Supplier\_name List\_price





## Summary

- Mapping from ERDs to relations is an algorithmic process
- Some choice points involve comparing time-space tradeoffs (more in physical design)

