

Question2: Explain the difference between an I/O-bound process and a CPU-bound process. 2 points

Ans: The differences between the two types of processes stem from the number of I/O requests that the process generates. An I/O-bound process spends more of its time seeking I/O operations than doing computational work. The CPU-bound process infrequently requests I/O operations and spends more of its time performing computational work.

Question3: Describe how UNIX and Linux manage orphan processes. 2 points

Ans: If a parent terminates without first calling wait(), its children are considered orphan processes. Linux and UNIX assign the init process as the new parent of orphan processes and init periodically calls wait() which allows any resources allocated to terminated processes to be reclaimed by the operating system.

Question4: What are the three general ways that a deadlock can be handled? 3 points

Ans:

- 1- A deadlock can be prevented by using protocols to ensure that a deadlock will never occur.
- 2- A system may allow a deadlock to occur, detect it, and recover from it.
- 3- operating system may just ignore the problem and pretend that deadlocks can never occur.

Question 5: Consider the following snapshot of a system: 4 points

	Allocation	Max	Available
	<i>A B C D</i>	<i>A B C D</i>	<i>A B C D</i>
<i>P0</i>	2 0 0 1	4 2 1 2	3 3 2 1
<i>P1</i>	3 1 2 1	5 2 5 2	
<i>P2</i>	2 1 0 3	2 3 1 6	
<i>P3</i>	1 3 1 2	1 4 2 4	
<i>P4</i>	1 4 3 2	3 6 6 5	

If a request from process *P1* arrives for (1, 1, 0, 0), can the request be granted immediately?

Allocate *p1* the required resources

	Allocation	Max	Available
	<i>A B C D</i>	<i>A B C D</i>	<i>A B C D</i>
<i>P0</i>	2 0 0 1	4 2 1 2	2 2 2 1
<i>P1</i>	4 2 2 1	5 2 5 2	
<i>P2</i>	2 1 0 3	2 3 1 6	
<i>P3</i>	1 3 1 2	1 4 2 4	
<i>P4</i>	1 4 3 2	3 6 6 5	

Apply banker algorithm on the matrix

	NEED	work	
	<i>A B C D</i>	2 2 2 1	
<i>P0</i>	2 2 1 1		1- process P0 can be served then marked as finished and work become 4 2 2 1
<i>P1</i>	1 0 3 1		2- process P1 can be served then marked as finished and work become 8 4 4 2
<i>P2</i>	0 2 1 3		3- process P2 can be served then marked as finished and work become 10 5 4 3
<i>P3</i>	0 1 1 2		4- process P3 can be served then marked as finished and work become 11 8 5 5
<i>P4</i>	2 2 3 3		5- process P4 can be served then marked as finished and work become 12 12 8 7

It is safe to allocate *p1* (1,1,0,0) so the request can be granted immediately.

Question 6:How is a limit register used for protecting main memory?

2 points

Ans: When the CPU is executing a process, it generates a logical memory address that is added to a relocation register in order to arrive at the physical memory address actually used by main memory. A limit register holds the maximum logical address that the CPU should be able to access. If any logical address is greater than or equal to the value in the limit register, then the logical address is a dangerous address and an error results.

Question 7:Explain how enhanced second chance algorithm is works.

4 points

The algorithm define reference bit and the modify bit (described in Section 9.4.1) as an ordered pair. With these two bits, we have the following four possible classes:

1. (0, 0) neither recently used nor modified—best page to replace
2. (0, 1) not recently used but modified—not quite as good, because the page will need to be written out before
3. (1, 0) recently used but clean—probably will be used again soon
4. (1, 1) recently used and modified—probably will be used again soon, and the page will be need to be written out to disk before it can be replaced

Question 8:Consider the following page reference string: 1, 2, 3, 4, 2, 1, 5, 6, 2, 1, 2, 3, 7, 6, 3, 2, 1, 2, 3, 6. How many page faults would occur for the following replacement algorithms, Remember that all frames are initially empty.

6 points

• LRU replacement

1	1	1	4	4	4	5	5	5	1	1	1	1	7	7	7	2	2	2	2
	2	2	2	2	2	2	6	6	6	6	6	3	3	3	3	3	3	3	3
		3	3	3	1	1	1	2	2	2	2	2	2	6	6	6	1	1	6
PF	PF	PF	PF		PF	PF	PF	PF	PF			PF	PF	PF		PF	PF		PF

Number of page fault = 15

• FIFO replacement

1	1	1	4	4	4	4	6	6	6	6	3	3	3	3	2	2	2	2	6
	2	2	2	2	1	1	1	2	2	2	2	7	7	7	7	1	1	1	1
		3	3	3	3	5	5	5	1	1	1	1	6	6	6	6	6	3	3
PF	PF	PF	PF		PF	PF	PF	PF	PF		PF	PF	PF		PF	PF		PF	PF

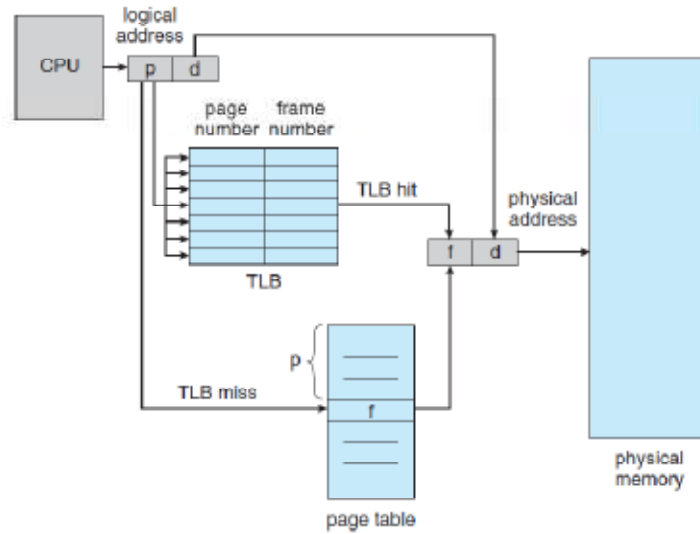
Number of page fault = 16

• Optimal replacement

1	1	1	1	1	1	1	1	1	1	1	3	3	3	3	3	3	3	3	6
	2	2	2	2	2	2	2	2	2	2	2	7	7	7	2	2	2	2	2
		3	4	4	4	5	6	6	6	6	6	6	6	6	6	1	1	1	1
PF	PF	PF	PF			PF	PF				PF	PF			PF	PF			PF

Number of page fault = 11

Question 9: Describe the address translation process, show how TLB can be used with page table to find the physical address from the logical address. 3 points



Question 10: How files can be protected from race condition when multiple access occur. 2 points

Ans:

By using shared lock (which is a read lock) and exclusive Lock (which is write lock)
 No update to the file while shared lock is acquired by a process or more and no read or write on the file while the exclusive lock is acquired by any process

Question 11: Given the following code write a code for test_and_set function below 2 points

```
do {
while (test_and_set(&lock));

/* critical section */

lock = false;

} while (true);
```

Ans.

```
boolean test and set(boolean *target) {
boolean rv = *target;
*target = true;
return rv;
}
```